

Performance Optimization: Simulation and Real Measurement

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Agenda

- Introduction
- Performance Analysis
- Profiling Tools: Examples & Demo
- KCachegrind: Visualizing Results
- What's to come ...



Introduction

- Why Performance Analysis in KDE ?
 - Key to useful Optimizations
 - Responsive Applications required for Acceptance
 - Not everybody owns a P4 @ 3 GHz
- About Me
 - Supporter of KDE since Beginning (“KAbalone”)
 - Currently at TU Munich, working on Cache Optimization for Numerical Code & Tools



Agenda

- Introduction
- Performance Analysis
 - Basics, Terms and Methods
 - Hardware Support
- Profiling Tools: Examples & Demo
- KCachegrind: Visualizing Results
- What's to come ...



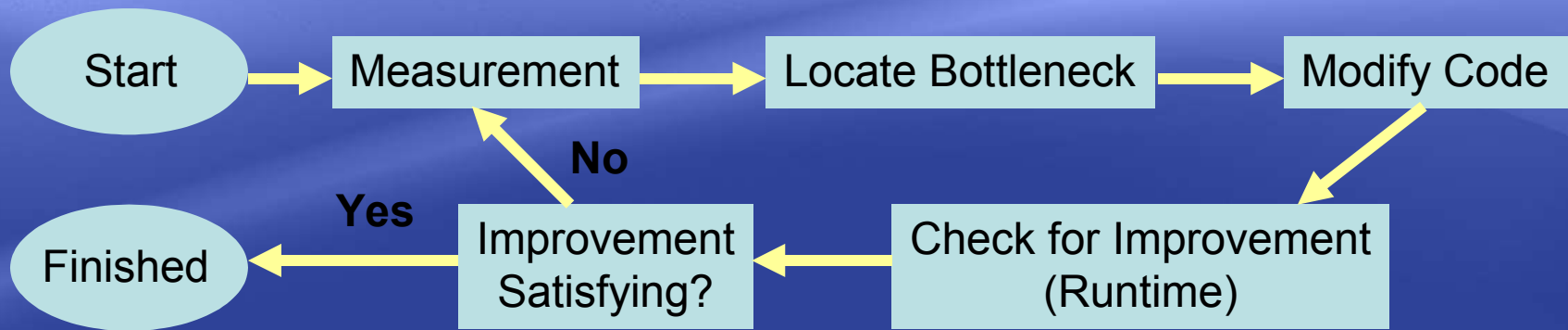
Performance Analysis

- Why to use...
 - Locate Code Regions for Optimizations (Calls to time-intensive Library-Functions)
 - Check for Assumptions on Runtime Behavior (same Paint-Operation multiple times?)
 - Best Algorithm from Alternatives for a given Problem
 - Get Knowledge about unknown Code (includes used Libraries like KDE-Libs/QT)



Performance Analysis (Cont'd)

- How to do...
 - At End of (fully tested) Implementation
 - On Compiler-Optimized Release Version
 - With typical/representative Input Data
 - Steps of Optimization Cycle



Performance Analysis (Cont'd)

- Performance Bottlenecks (sequential)
 - Logical Errors: Too often called Functions
 - Algorithms with bad Complexity or Implementation
 - Bad Memory Access Behavior (Bad Layout, Low Locality) ← Too low-level for GUI Applications ?
 - Lots of (conditional) Jumps, Lots of (unnecessary) Data Dependencies, ... ↙



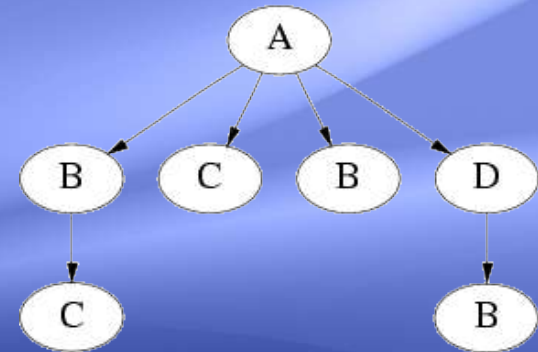
Performance Measurement

- Wanted:
 - Time Partitioning with
 - Reason for Performance Loss (Stall because of...)
 - Detailed Relation to Source (Code, Data Structure)
 - Runtime Numbers
 - Call Relationships, Call Numbers
 - Loop Iterations, Jump Counts
 - No Perturbation of Results b/o Measurement



Measurement - Terms

- Trace: Stream of Time-Stamped Events
 - Enter/Leave of Code Region, Actions, ...
- Example: Dynamic Call Tree



- Huge Amount of Data (Linear to Runtime)
- Unneeded for Sequential Analysis (?)



Measurement – Terms (Cont'd)

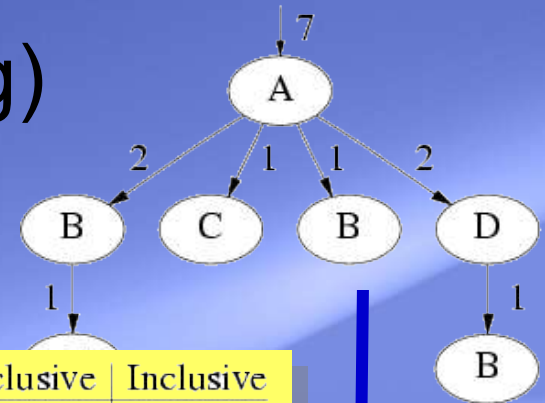
- Profiling (e.g. Time Partitioning)

- Summary over Execution

- Exclusive, Inclusive Cost / Time, Counters

- Example:
DCT → DCG
(Dynamic Call Graph)

- Amount of Data
Linear to Code Size



	Exclusive	Inclusive
A	1	7
D	1	2
B	3	4
C	2	2



Methods

- Precise Measurements
 - Increment Counter (Array) on Event
 - Attribute Counters to
 - Code / Data
 - Data Reduction Possibilities
 - Selection (Event Type, Code/Data Range)
 - Online Processing (Compression, ...)
 - Needs Instrumentation (Measurement Code)



Methods - Instrumentation

- Manual
- Source Instrumentation
- Library Version with Instrumentation
- Compiler
- Binary Editing
- Runtime Instrumentation / Compiler
- Runtime Injection



Methods (Cont'd)

- Statistical Measurement (“Sampling”)
 - TBS (Time Based), EBS (Event Based)
 - Assumption: Event Distribution over Code
Approximated by checking every N-th Event
 - Similar Way for Iterative Code:
Measure only every N-th Iteration
- Data Reduction Tunable
 - Compromise between Quality/Overhead



Methods (Cont'd)

- Simulation
 - Events for (not existant) HW Models
 - Results not influenced by Measurement
 - Compromise Quality / Slowdown
 - Rough Model = High Discrepancy to Reality
 - Detailed Model = Best Match to Reality
But: Reality (CPU) often unknown...
 - Allows for Architecture Parameter Studies



Hardware Support

- Monitor Hardware
 - Event Sensors (in CPU, on Board)
 - Event Processing / Collection / Storing
 - Best: Separate HW
 - Comprimise: Use Same Resources after Data Reduction
 - Most CPUs nowadays include Performance Counters



Performance Counters

- Multiple Event Sensors
 - ALU Utilization, Branch Prediction, Cache Events (L1/L2/TLB), Bus Utilization
- Processing Hardware
 - Counter Registers
 - Itanium2: 4, Pentium-4: 18, Opteron: 8
Athlon: 4, Pentium-II/III/M: 2, Alpha 21164: 3



Performance Counters (Cont'd)

- Two Uses:
 - Read
 - Get Precise Count of Events in Code Regions by Enter/Leave Instrumentation
 - Interrupt on Overflow
 - Allows Statistical Sampling
 - Handler Gets Process State & Restarts Counter
- Both can have Overhead
- Often Difficult to Understand



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- Profiling Tools: Examples & Demo
 - Callgrind/Calltree
 - OProfile
- KCachegrind: Visualizing Results
- What's to come ...



Tools - Measurement

- Read Hardware Performance Counters
 - Specific: **PerfCtr** (x86), Pfmon (Itanium), perfex (SGI)
 - Portable: PAPI, PCL
- Statistical Sampling
 - PAPI, Pfmon (Itanium), **OProfile** (Linux), VTune (commercial - Intel), Prof/GProf (TBS)
- Instrumentation
 - GProf, Pixie (HP/SGI), VTune (Intel)
 - DynaProf (Using DynInst), **Valgrind** (x86 Simulation)



Tools – Example 1

- GProf (Compiler generated Instr.):
 - Function Entries increment Call Counter for (caller, called)-Tupel
 - Combined with Time Based Sampling
 - Compile with “gcc -pg ...”
 - Run creates “gmon.out”
 - Analyse with “gprof ...”
 - Overhead still around 100% !
- Available with GCC on UNIX



Tools – Example 2

- Callgrind/Calltree (Linux/x86), GPL
 - Cache Simulator using Valgrind
 - Builds up Dynamic Call Graph
 - Comfortable Runtime Instrumentation
 - <http://kccachegrind.sf.net>
- Disadvantages
 - Time Estimation Inaccurate
(No Simulation of modern CPU Characteristics!)
 - Only User-Level



Tools – Example 2 (Cont'd)

- Callgrind/Calltree (Linux/x86), GPL
 - Run with “callgrind prog”
 - Generates “callgrind.out.xxx”
 - Results with “callgrind_annotate” or “kcachegrind”
 - Cope with Slowness of Simulation:
 - Switch of Cache Simulation: `--simulate-cache=no`
 - Use “Fast Forward”:
`--instr-atstart=no / callgrind_control -i on`
- DEMO: KHTML Rendering...



Tools – Example 3

- **OProfile**

- Configure (as Root: `oprof_start`, `~/.oprofile/daemonrc`)
- Start the OProfile daemon (`opcontrol -s`)
- Run your code
- Flush Measurement, Stop daemon (`opcontrol -d/-h`)
- Use tools to analyze the profiling data
 - `opreport`: Breakdown of CPU time by procedures (better: `opreport -gdf | op2calltree`)

- DEMO: KHTML Rendering...



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- KCachegrind: Visualizing Results
 - Data Model, GUI Elements, Basic Usage
 - DEMO
- What's to come ...



KCachegrind – Data Model

- Hierarchy of Cost Items (=Code Relations)
 - Profile Measurement Data
 - Profile Data Dumps
 - Function Groups:
Source files, Shared Libs, C++ classes
 - Functions
 - Source Lines
 - Assembler Instructions



KCachegrind – GUI Elements

- List of Functions / Function Groups
- Visualizations for an Activated Function

- DEMO

Event Type	Incl.	Self	Short	Formula
L1 Instr. Fetch Miss	49.26	1.83	l1mr	
L1 Data Read Miss	100.00	0.00	D1mr	
L1 Data Write Miss	99.99	0.00	D1mw	
L2 Instr. Fetch Miss	49.20	1.81	l2mr	
L2 Data Read Miss	100.00	0.00	D2mr	
L2 Data Write Miss	99.99	0.00	D2mw	
L1 Miss Sum	100.00	0.00		$L1m = l1mr + D1mr + D1mw$
L2 Miss Sum	100.00	0.00		$L2m = l2mr + D2mr + D2mw$
Cycle Estimation	100.00	0.00		$CEst = Ir + 10 L1m + 100 L2m$

#	CEst	Source ('mg.f')
378		C Get down
379	0.00	do k= l,1,-1
380		if((k.EQ.l).OR.(1.EQ.1).OR.(0.EQ.1)) then
381		C for smoothing function selection see config.h
382	0.00	call rb4w(k,2,wsp(uoff(k)),wsp(fcoff(k)),
	35.21	150 calls to 'rb4w_' (mg)
383	0.00	& (xdim(k)-1),(ydim(k)-1),(zdim(k)-1))
384		else
385		C for other grids use std (only for only_u runs!)
386		call rb_wf(k,2,wsp(uoff(k)),wsp(fcoff(k)),

#	CEst	Hex	Assembler
804 AC3D	0.00 56		push %esi
804 AC3E	0.00 57		push %edi
804 AC3F	0.00 50		push %eax
804 AC40	0.00 68 4c 20 0b 08		push \$0x80b204c
804 AC45	0.00 52		push %edx
804 AC46	0.00 e8 8b b2 00 00		call 8055ed6 <rb4w_
	35.21		150 calls to 'rb
804 AC4B	0.00 8d 44 24 1c		lea 0x1c(%esp),%eax
804 AC4F	0.00 50		push %eax



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- What's to come ...
 - Callgrind
 - KCachegrind



What's to come

- Callgrind
 - Free definable User Costs
("MyCost += arg1" on Entering MyFunc)
 - Relation of Events to Data Objects/Structures
 - More Optional Simulation (TLB, HW Prefetch)



What's to come (Cont'd)

- KCachegrind
 - Supplement Sampling Data with Inclusive Cost via Call-Graph from Simulation
 - Comparison of Measurements
 - Plugins for
 - Interactive Control of Profiling Tools
 - Visualizations
- Visualizations for Data Relation



Finally...

THANKS FOR LISTENING



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